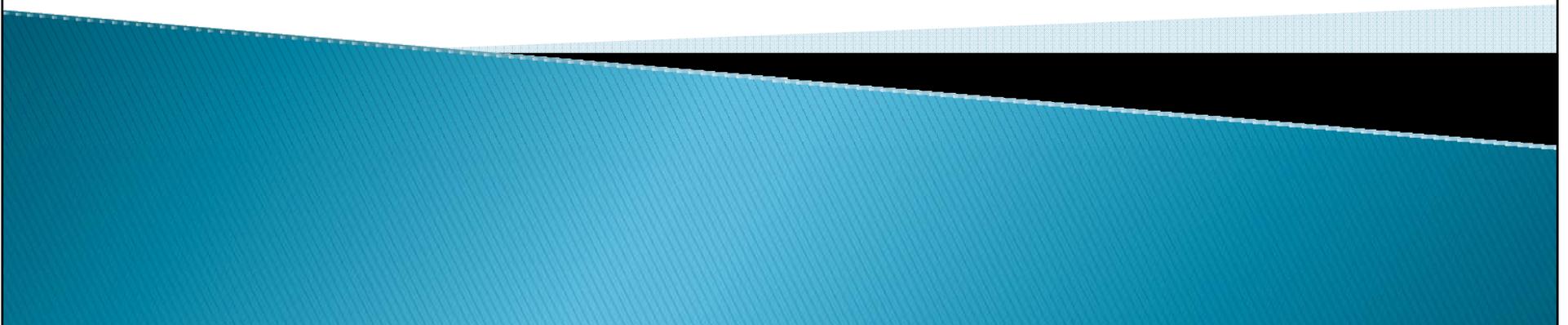
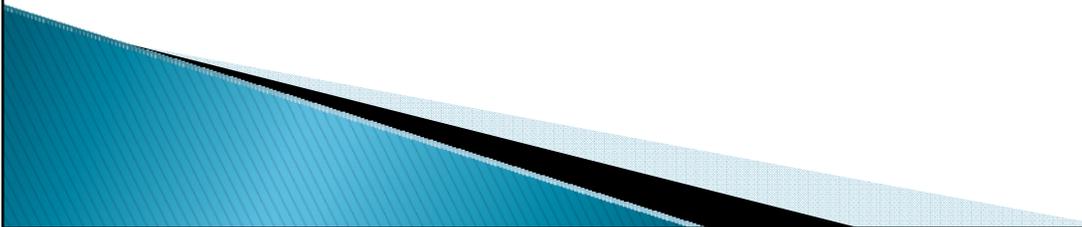


# CSSE 220 Day 14

More algorithm efficiency analysis, Big-Oh  
Work on Paint



# CSSE 220 Day 14

- ▶ **BallWorlds grades** are on ANGEL, and scoring sheets handed back.
  - ▶ **Another progress Report/IEP** due today at the end of class.
  - ▶ Today:
    - Algorithm analysis and Big-Oh
    - More Paint time
  - ▶ Questions?
- 

# Questions on Angel part of exam?

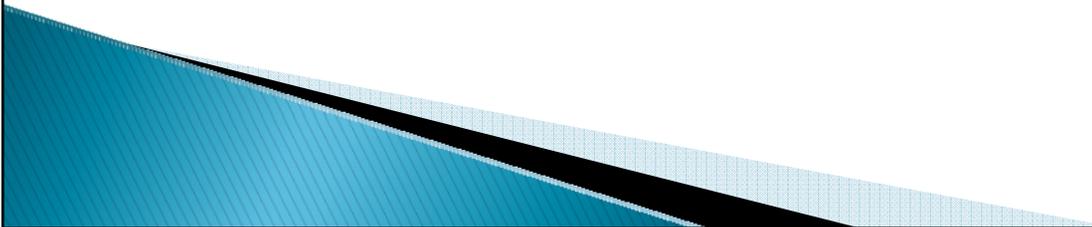
- ▶ Or on programming part?
- ▶ This is all fair game for the Final Exam, only 6 short weeks away (**Monday, 6pm in Crapo G310**).

# Key Concepts quiz tomorrow

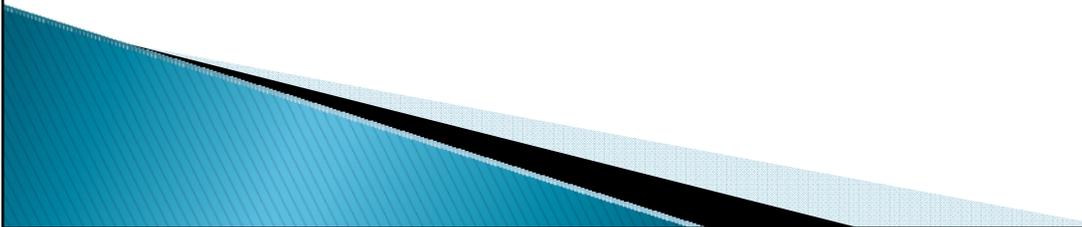
- ▶ It is important that we not only be able to write object-oriented programs, but that we build a **vocabulary** that enables us to communicate with each other about them.
- ▶ That is why we asked you to spend four weeks learning the "lingo" of OOP in Java.
- ▶ Tomorrow is the check-up on that.
- ▶ This ANGEL-based quiz is closed book and notes.
- ▶ It consists of matching questions, and you will only have about 30 seconds per term to complete it. So know your terms well!

# Measuring program efficiency

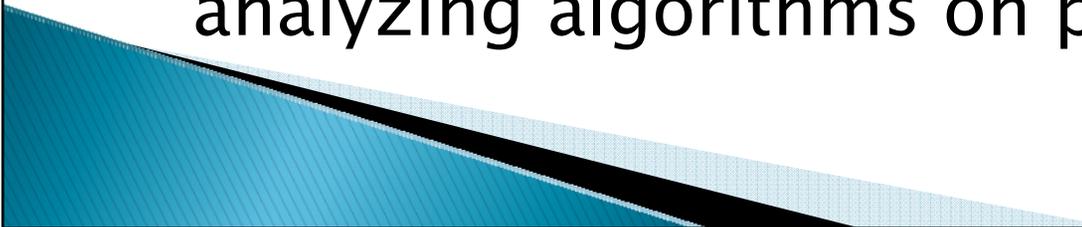
- ▶ What kinds of things should we measure?
  1. CPU time
  2. memory used
  3. disk transfers
  4. network bandwidth
- ▶ Mostly in this course, we focus on the first two, and especially on CPU time.



# Some simple efficiency tips

- If a statement in a loop calculates the same value each time through, move it outside the loop
  - Store and retain data on a “need to know” basis.
  - Don't store what you won't reuse!
    - Do store what you need to reuse!
  - Don't put everything into an array when you only need one or two consecutive items at a time.
  - Don't make a variable be a field when it can be a local variable of a method.
- 

# An example of running time

- ▶ How can we measure running time?
  - ▶ `System.currentTimeMillis`
  - ▶ Run Sieve example.
  - ▶ When do we really care about efficiency?
  - ▶ Can we get a rough idea of efficiency by analyzing algorithms on paper?
- 

## Familiar example:

Linear search of a sorted array of Comparable items

```
for (int i=0; i < a.length; i++)
    if ( a[i].compareTo(soughtItem) > 0 )
        return NOT_FOUND;
    else if ( a[i].compareTo(soughtItem) == 0 )
        return i;
return NOT_FOUND;
```

- What should we count?
- Best case, worst case, average case?

# Another algorithm analysis example

Does the following method actually create and return a copy of the string *s*?

What can we say about the running time of the method?  
(where *N* is the length of the string *s*)

**What should we count?**

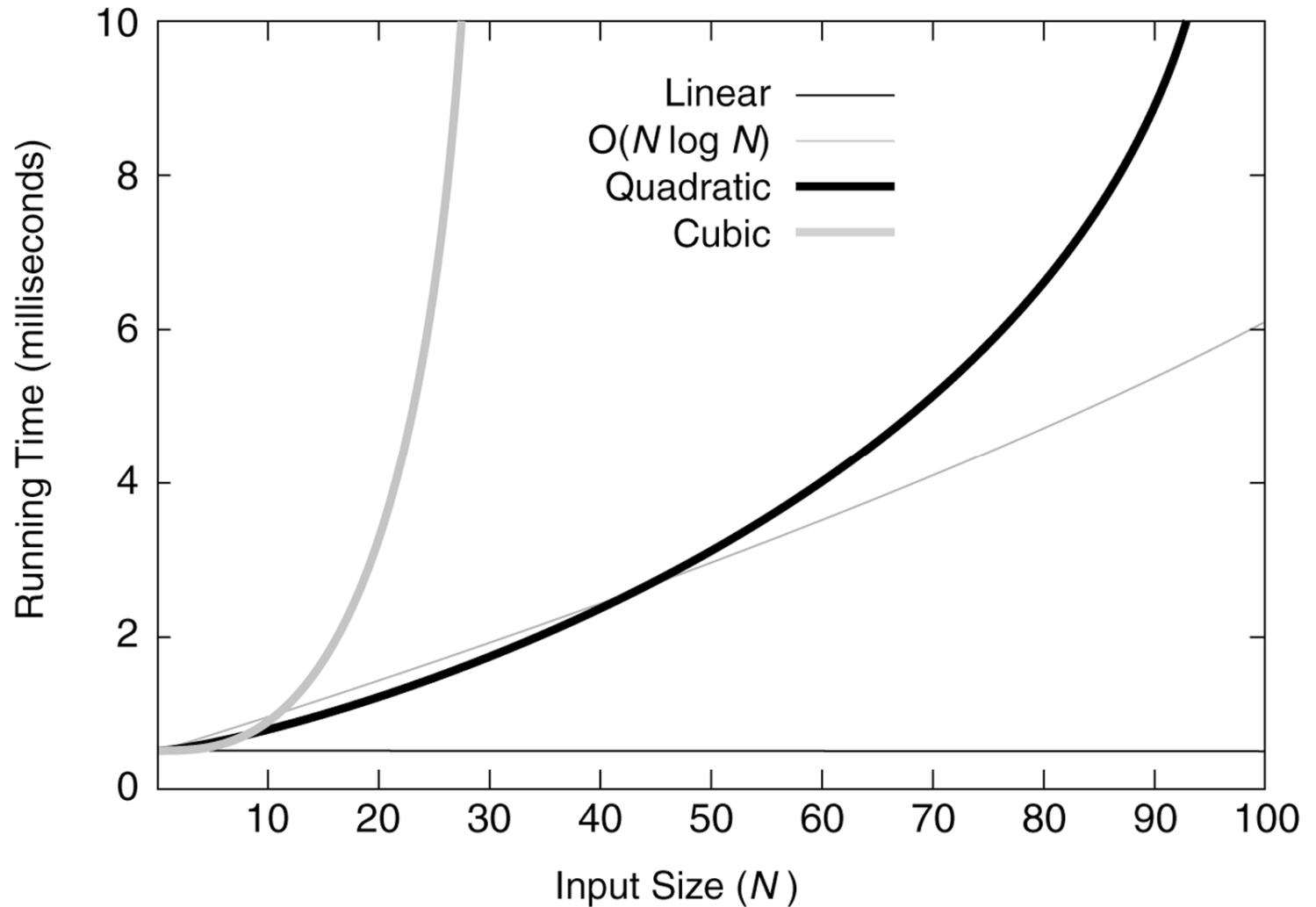
```
public static String stringCopy(String s) {  
    String result = "";  
    for (int i=0; i<s.length(); i++)  
        result += s.charAt(i);  
    return result;  
}
```

Don't be too quick to make assumptions when  
analyzing an algorithm!

**How can we do the copy more efficiently?**

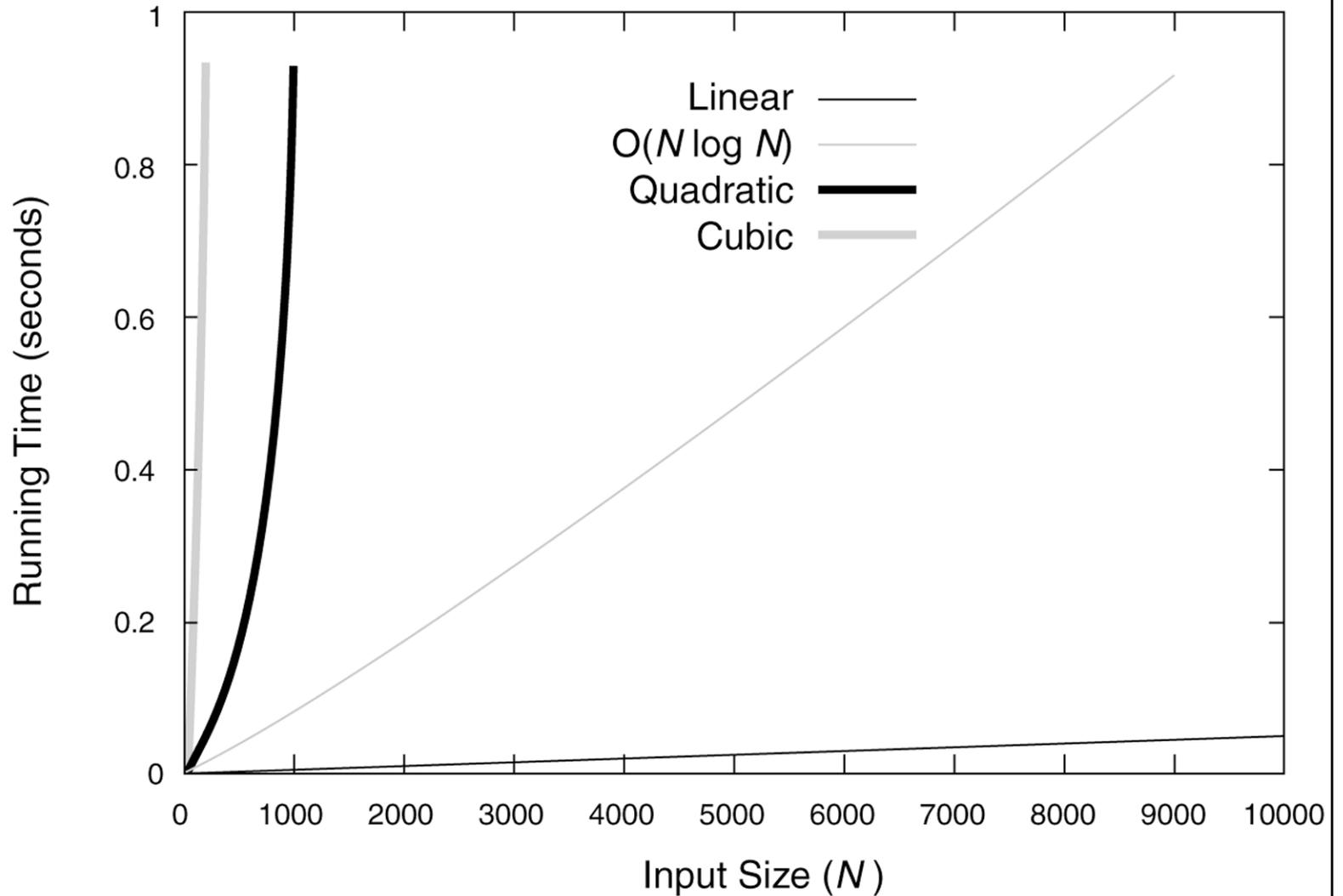
## Figure 5.1

Running times for small inputs



## Figure 5.2

Running times for moderate inputs



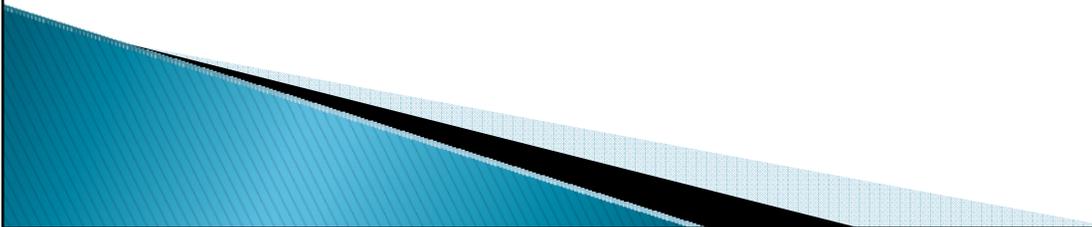
## Figure 5.3

Functions in order of increasing growth rate

FUNCTION	NAME
$c$	Constant
$\log N$	Logarithmic
$\log^2 N$	Log-squared
$N$	Linear
$N \log N$	$N \log N$ ← a.k.a "log linear"
$N^2$	Quadratic
$N^3$	Cubic
$2^N$	Exponential

# Asymptotic analysis

- ▶ We only really care what happens when  $N$  (the size of a problem) gets large.
- ▶ Is the function linear? quadratic? etc.

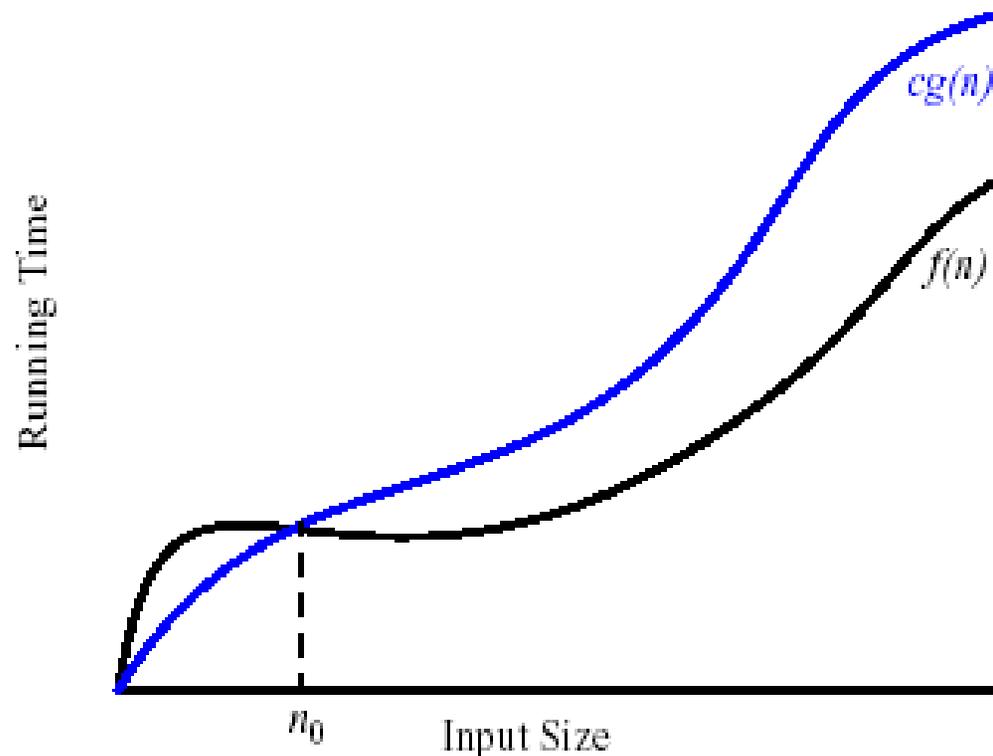


- The “Big-Oh” Notation

- given functions  $f(n)$  and  $g(n)$ , we say that  $f(n)$  is  $O(g(n))$  if and only if  $f(n) \leq c g(n)$  for  $n \geq n_0$

- $c$  and  $n_0$  are constants,  $f(n)$  and  $g(n)$  are functions over non-negative integers

In this course, we won't be so formal . We'll just say that  $f(N)$  is  $O(g(N))$  means that  $f(n)$  is eventually smaller than a constant times  $g(n)$ .



- **Simple Rule:** Drop lower order terms and constant factors.
  - $7n - 3$  is  $O(n)$
  - $8n^2 \log n + 5n^2 + n$  is  $O(n^2 \log n)$
- Special classes of algorithms:
  - logarithmic:  $O(\log n)$
  - linear:  $O(n)$
  - quadratic:  $O(n^2)$
  - polynomial:  $O(n^k), k \geq 1$
  - exponential:  $O(a^n), n > 1$
- “Relatives” of the Big-Oh
  - $\Omega(f(n))$ : Big Omega
  - $\Theta(f(n))$ : Big Theta

# Limits and asymptotics

- ▶ consider the limit

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)}$$

- ▶ What does it say about asymptotics if this limit is zero, nonzero, infinite?
- ▶ We could say that knowing the limit is a sufficient but not necessary condition for recognizing big-oh relationships.
- ▶ It will be sufficient for all examples in this course.

# Apply this limit property to the following pairs of functions

Assume  $a$  and  $b$  are constants.

- $N$  and  $N^2$
- $N^2 + 3N + 2$  and  $N^2$
- $N + \sin(N)$  and  $N$
- $\log N$  and  $N$
- $N \log N$  and  $N^2$
- $N^a$  and  $N^n$
- $a^N$  and  $b^N$  ( $a < b$ )
- $\log_a N$  and  $\log_b N$  ( $a < b$ )
- $N!$  and  $N^N$

# Big-Oh Style

## ▶ Give tightest bound you can

- Saying that  $3N+2$  is  $O(N^3)$  is true, but not as useful as saying it's  $O(N)$  [What about  $\Theta(N^3)$  ?]

## ▶ Simplify:

- You *could* say:
- $3n+2$  is  $O(5n-3\log(n) + 17)$
- and it would be technically correct...
- It would also be poor taste ... and put me in a bad mood.

## ▶ But... if I ask “true or false: $3n+2$ is $O(n^3)$ ”, what’s the answer?

- True!
- There may be “trick” questions like this on assignments and exams.
- But they aren’t really tricks, just following the big-Oh definition!

# Rest of Class

- ▶ Please provide me with some mid-term feedback on how the course is going
  - Angel > Lessons > Course Discussion Forums > Midterm Plus/Delta Feedback
- ▶ Thanks!
- ▶ After that you have all class to work on Paint with your partner.
- ▶ Please commit your updated IEP at the end of class.